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Brook: An Easy and Efficient Framework for Distributed Machine Learning

Chao Ma, Yan Ni, and Zhen Xiao*
Department of Computer Science and Technology
Peking University
{machao, niyan, xiaozhen}@net.pku.edu.cn

Abstract

We present Brook, a new framework for distributed machine learning problems. Like some previous frameworks, Brook adopts the parameter server paradigm that simplifies the task of distributed programming. Unlike these frameworks, we build a novel system component called parameter agent that masks the communication details between workers and servers by mapping remote servers to local in-memory file. In this way, Brook provides a simple and platform-independent interface called RWW, where users can migrate existing single-machine programs, written in any programming language, to the distributed environment with minimal modification. In addition, to achieve system efficiency and scalability, Brook is designed to minimize network traffic, maximize CPU and memory utilizations, and support flexible fault-tolerance strategies. Our evaluation results show that Brook has the highly competitive performance and scalability, while providing enhanced ease of use compared to existing frameworks.

1 Introduction

Machine learning (ML) is becoming the primary mechanism to extract useful knowledge from Big Data. To improve the accuracy, ML methods tend to use models with more parameters trained on large numbers of examples. However, due to the computation and storage limit of a single machine, executed in a distributed manner has become a prerequisite for solving large-scale ML problems.

The data-flow frameworks such as Hadoop [4] and Spark [2], have significantly simplified the task of building large-scale data processing on commodity clusters. Based on these frameworks, the distributed ML libraries such as Mahout [5] and MLI [6], have been widely used in both academia and industry. However, most of these frameworks adopt the iterative MapReduce [1] paradigm that mandates synchronous and coarse-grained computation and communication. This inherent system design for batch tasks incurs great inefficiency and is often inadequate when building the “Big Model” ML applications such as large-scale sparse logistic regression, massive topic model and deep networks. Parameter server paradigm [7] has recently emerged as an efficient approach to resolve the “Big Model” ML challenge. Under this paradigm, both the training data and workloads are spread across worker nodes, while the server nodes maintain the globally shared parameters. In contrast to the iterative MapReduce paradigm, computation and communication in parameter server can be asynchronous and fine-grained, and hence can improve the CPU utilization and reduce the communication cost dramatically.

The aforementioned frameworks have proven to be tremendously useful for simplifying the task of building distributed ML applications. However, almost all of them force users to re-write their existing code in a new software stack, many of which expose an unfamiliar programming environment to

*The contact author is Zhen Xiao.

054 users. For example, on one hand, many ML developers are accustomed to using powerful, high pro-
 055 ductivity array-languages such as Matlab, R and Numpy. For these users, especially inexperienced
 056 ones, the steep learning curve of the new programming platforms, including both the programming
 057 language and the programming model, is a major obstacle to the adoption of the new frameworks.
 058 On the other hand, some skilled developers in ML domain prefer to use more efficient programming
 059 languages such as C/C++, as well as the high-performance hardwares such as GPGPU to build ML
 060 applications which can be tuned to give extremely good performance. For these users, most of the
 061 popular frameworks, such as Spark, cannot satisfy their needs. So here comes a natural question:
 062 can we build a new framework where users can re-use their single-machine code easily and the
 063 framework is agnostic of the underlying programming platforms?

064 To answer this question, in this paper we propose Brook, a new framework which allows developers
 065 to completely get rid of the restrictions of the programming platforms. Using Brook, developers can
 066 migrate existing single-machine ML programs, potentially written in any programming language
 067 or executed on the specific underlying hardwares such as GPGPU, to a distributed environment for
 068 concurrent execution with little change, while achieving the similar fault-tolerance guarantees and
 069 the enhanced performance compared to existing frameworks. To achieve these goals, Brook extends
 070 the original parameter server paradigm by adding a novel component called parameter agent and its
 071 counterparts. Section 2 explains how these components work closely together and demonstrates how
 072 users build ML applications via the simple and language-agnostic interface RWW. Furthermore, to
 073 achieve high performance and scalability, both computational intensive workloads and the volume
 074 of data communication demand careful system design and optimization. In section 3, we discuss
 075 the optimization approaches used in Brook, including the vector store, message compression, and
 076 flexible fault-tolerance strategies. In the last section, we show a preliminary performance evaluation
 077 of our system.

078 2 System Architecture

080 This section describes our system architecture. In our presentation, we first provide an overview
 081 of Brook (§2.1) and then introduce the programming model based on the RWW interface (§2.2).
 082 After that, we make a comparison between Brook and the existing solutions for the cross-language
 083 programming in ML systems, and show the advantages of Brook (§2.3). Finally, we discuss the
 084 implementation of our system (§2.4).

086 2.1 Overview

088 Brook’s execution environment consists of one
 089 master process and many server and worker
 090 processes, each executing on a potentially dif-
 091 ferent machine. Figure 1 illustrates the over-
 092 all interactions among the master, servers and
 093 workers when executing a Brook program. As
 094 Figure 1 shows, a server node stores assigned
 095 parameters partition in its memory and handles
 096 the aggregation and update operations associ-
 097 ated with that partition. Each worker node is
 098 responsible for storing a portion of the training
 099 data to compute local updates such as gradient.
 100 The master node maintains the bookkeeping of
 101 each worker and server process, which can be
 102 recovered without interrupting the computation
 103 when it crashes by non-catastrophic machine
 104 failures. Similar to existing frameworks, we assume that the master node failures are rare and hence
 105 provide no protection for that. Note that workers communicate only with the servers and the master,
 106 not among themselves.

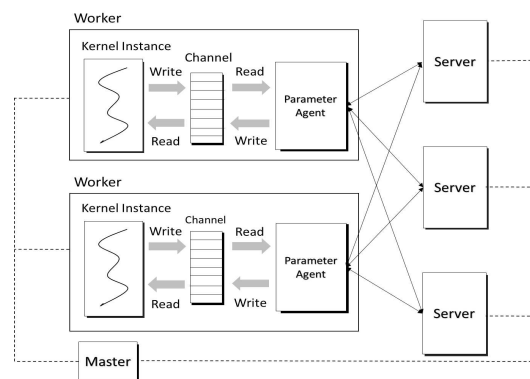


Figure 1: Brook Architecture

106 In Figure 1, a ML task is divided among all of the workers, which jointly learn the globally shared
 107 parameters by computing the local updates based on its own training data. These local updates can
 then be sent to servers via a specific partition algorithm in the worker nodes. Servers aggregate all

108 of these updates before applying them to the corresponding parameters and later send these updated
 109 parameters back to the corresponding workers, as the response of their requests. This series of
 110 processes is similar to that in the original parameter server.

111 The main difference between Brook and the original parameter server model is that the client process
 112 in Brook dose not communicate with the server nodes directly. Instead, at run-time each worker
 113 process derives two child processes, including the client (also called kernel instance in Brook) and
 114 the parameter agent. Workers communicate with the servers through the parameter agent processes
 115 and Brook combines the kernel instance and the parameter agent with the communication channel,
 116 which is a platform-independent abstraction for local data transport and synchronization. We now
 117 describe the details of these components in terms of the basic functions and the interactions of them.

118 **Kernel Instance.** Kernel instance is provided by the application developer and started as an inde-
 119 pendent process that can be executed on any language platform. The effect of a kernel instance is to
 120 repeatedly read new parameters from channel, compute local updates, and write both the updates
 121 and the request to channel. As will be discussed later, because the parameter agent abstraction hides the
 122 communication details, the distributed shared parameters appear to be local. This allows the kernel
 123 instance to easily retro-fit existing single-machine implementations with minimal modification.

124 **Parameter Agent.** Parameter agent acts as a middleman between the client and the servers. Brook
 125 makes use of the abstraction of the parameter agent to simplify the logic of the client process. The
 126 cumbersome system work such as network communication, message queue and serialization will
 127 be taken over by the parameter agent process. By doing so, a client process can focus only on the
 128 implementation of the core algorithm in ML, based on any underlying programming platform, and
 129 just read from and write to channel via the simple and language-agnostic interface to share data with
 130 the parameter agent process.

131 **Channel.** Channel consists of a data channel which is responsible for local data transport, and a
 132 signal channel which manages the synchronization. The concrete implementation of the channel
 133 abstraction is file, since all programming languages can access it in a consistent way. There are two
 134 types of files used in a channel. For the data channel, we use the in-memory file which is like a
 135 regular file but based on Ramfs with at least two orders of magnitude improvement in performance.
 136 For the signal channel, we use FIFO, which is easy to perform blocking I/O between two processes.

137 Based on the foregoing descriptions, in the next section we demonstrate how users build their dis-
 138 tributed ML applications on Brook by using the RWW interface.

140 2.2 Programming model

142 Although ML algorithms come in many forms,
 143 almost all of them seek a set of parameters to
 144 a global model A that best summarizes or ex-
 145 plains the input data \mathcal{D} . Such problems are usu-
 146 ally solved by iterative-convergent algorithms,
 147 many of which can be abstracted as the addi-
 148 tive form: $A^{(t+1)} = A^{(t)} + \Delta(A^{(t)}, \mathcal{D})$, where
 149 $A^{(t)}$ is the state of model parameters at iteration
 150 t , \mathcal{D} is the input data, and the kernel function
 151 Δ computes the model updates from \mathcal{D} . Using
 152 Brook, developers focus on the implementation
 153 of the kernel instance, where the program will
 154 be distributed over worker nodes and run con-
 155 currently at run-time for computing their local
 156 updates such as gradients. As shown in Figure 2,
 157 we see that developers can easily migrate exist-
 158 ing single-machine code to the Brook environ-
 159 ment by invoking a set of functions called RWW
 160 at the end of each iteration in a ML task. These
 161 functions include the Write, the Read and the Wait.
 As we mentioned before, the RWW interface gives
 developers an easy and platform-independent way
 to communicate with the parameter agent process.
 Figure 2 illustrates that, the Write and the Read
 functions are responsible for local data transport
 by writing to and reading from data channel, while
 the Wait function manages the process synchroniza-
 tion through the signal channel by getting and
 setting the iteration timestamps (also called vec-
 tor clock) during Brook’s run-time. We note that, the

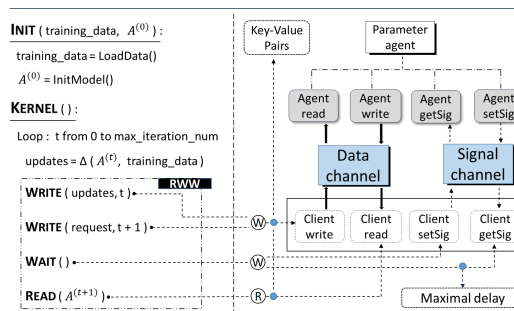


Figure 2: Programming model

162 RWW interface could be implemented either by users or by system developers, since the protocol of
163 data transport and synchronism that used in this interface are extremely simple.

164 Several features about the RWW interface are also worth noting. First, Brook supports a number of
165 data types for the cross-platform data transport, such as the text data, the 3rd-party message libraries
166 such as protobuf and Thrift, as well as the native binary data. Second, developers can configure
167 the value of the maximal delay in the Wait function, which could give developers the opportunity
168 to implement different consistency models such as the BSP, the SSP [24] and the ASP. The core
169 technique under the flexible consistency model mechanism is a signal queue. Furthermore, we
170 support user-defined update mechanism on server nodes by using the expression template, a simple
171 and efficient programming language trick of C++.

172 173 **2.3 Related solutions**

174 In contrast to our system, there are two common ways of the cross-language programming that
175 have been used in existing ML frameworks. The first way, many ML frameworks make use of the
176 languages wrapper around their original API. Such as a python wrapper for the original C++ API
177 by using the Boost or the SWIG [9] library. In contrast to Brook, the disadvantages of this solution
178 are obvious. First, building the languages wrapper is not an easy job, since we have to modify the
179 source code of the original framework and it might be inadequate for the inexperienced users who
180 need to do this by themselves. A real example in the open source community is the SparkR [3],
181 which is published in the latest version of Spark recently. So, in other words, it is unrealistic to build
182 the wrappers for every programming languages which might be used by existing single-machine ML
183 applications. However in Brook, the RWW interface is totally language-agnostic. Apart from this,
184 building languages wrapper naively is often lack of both efficiency and flexibility.

185 The second way, that is, the Hadoop streaming [10], which is widely used as it frees programmers
186 from Java language, which makes developers use power of Hadoop more easily. Similar to Brook,
187 Hadoop streaming uses the 3rd-party medium (standard input/output) to transport data over pro-
188 cesses. However, as we mentioned before, the inherent system design of Hadoop is not suitable for
189 the ML applications in both the programming model and the system performance. In addition, Hadoop
190 streaming transports all of the training data through the standard IO, which can incurs great system
191 overhead. By contrast, using Brook, we only transport the data of parameters through the channel.

192 193 **2.4 Implementation**

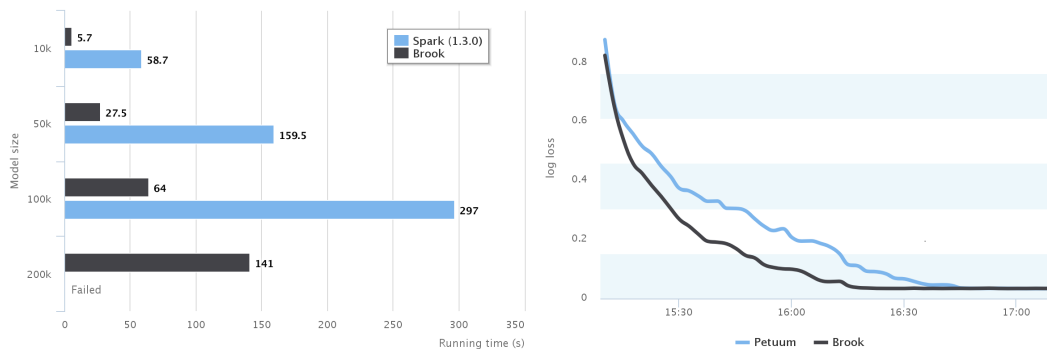
194 Brook is implemented in C++ and requires no change to the underlying OS or compiler. Our imple-
195 mentation re-uses a number of existing libraries such as MPICH2 for communication (supports both
196 the Ethernet and the InfiniBand network), Google's protobuf for object serialization, and Snappy
197 for data compression. Brook can run over the cluster resource manager such as Yarn [11] or Mesos
198 [12], and also provides an easy way for deployment on clusters by using the docker [13] container.

200 201 **3 System Optimization**

202 Implementing an efficient and scalable distributed framework is not an easy job, because both the
203 volume of data communication and the intensive computational workloads demand careful system
204 design and optimization. In this paper, we focus on three optimization techniques:

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206 **Message compression.** Since distributed machine learning problems typically require high band-
207 width, message compression is desirable. We use several compression approaches in our system to
208 reduce the network traffic as much as possible. First, we avoid sending single items because both
209 the communication overhead that caused by TCP/IP package header and the serialization overhead
210 are horrible. Hence we pack all of the single items into a batched message form. Second, instead of
211 using (key, value) pair to represent each item, the message consists of a list of (start-key, value-list),
212 where the value-list is a sequence of consecutive values. This simple modification can greatly reduce
213 the size of each message especially when message tends to be dense. Next, since many ML prob-
214 lems may use the same training data in different iterations, we cache the key lists in the receiving
215 nodes. Later the senders can send only a hash value of this list rather than the list itself. Finally, we
use the Snappy compression library to compress the serialized message.

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Figure 3: **Left:** System performance under the BSP consistency, versus Spark (1.3.0 version, 100 iterations), on logistic regression. **Right:** Convergence rate under the SSP consistency, versus Petuum (800-million parameters, maximal delay = 4), on large-scale sparse logistic regression. All datasets come from the Criteo CTR [19]. We note that, the left diagram shows that, the Spark task failed when we scale-up the mode size to 200k.

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Vector store. Many previous systems use key-value table to store the globally shared state during the run-time [23]. However, using this abstraction naively is inefficient in both the memory utilization and the computation. We found that operations in server nodes are typically represented as linear algebra computations. Hence in Brook, we use the vector store where the underlying implementation is the contiguous region of memory that stores only the consecutive values ordered by its potential index, while the non-existing items are associated with zeros automatically. Thus, we can save at least half of the memory cost and achieve efficiency by leveraging the high-performance multithreaded linear algebra libraries such as the OpenBLAS [15]. It also simplifies the system design for user-defined update mechanism on the server nodes. A similar approach is also used in the previous work [8].

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Flexible level of fault-tolerance. At scale, fault-tolerance is critical. However, most of the existing systems can only support one fixed fault-tolerance strategy, which is inflexible and often incurs much unnecessary overhead when system is deployed in a small, well-controlled cluster. Fortunately, in Brook, we give developers the opportunity to configure the fault-tolerance level, which could range from L0 to L3 and covers the cluster types from the small platform such as desktop PCs and lab-clusters, to the big, less predicabile platform such as large data center or the cloud. Brook will choose the different backup strategies such as backup on remote nodes or backup locally, according to the different fault-tolerance level.

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We will demonstrate at the workshop the concrete method about the flexible fault-tolerance strategies. We also plan to demo some other optimizations such as skip-list buffer, memory zero-copy and signal queue. Each of them has improved our system performance considerably.

255 4 Preliminary Evaluation

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We evaluated Brook on a cluster of 15 machines in our own laboratory. Each machine has a 8 cores Intel Xeon E5620 (2.40GHz) processor with 26GB memory. We compared Brook with Spark and Petuum under the BSP and the SSP consistency models, respectively. Our evaluation results show that Brook can outperform its competitors in both system efficiency and scalability. The highlights of our results can be found in Figure 3. During the workshop, we plan to demo more detailed evaluation of our system.

263 264 265 5 Acknowledgments

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